## **Goblint & GobPie** Towards Usable Data Race Verification

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#### Multiple Access data race Two threads simultaneously access same memory location...

$$T_1: lock(\&l_1);$$

v = v + 1;unlock(&l<sub>1</sub>);

$$T_{2}: lock(\&l_{1}); \\ v = v + 1; \\ unlock(\&l_{1}); \end{cases}$$

- List of accesses:  $\langle v, \{l_1\}, write, file.c: 2 \rangle$  $\langle v, \{l_1\}, write, file.c: 5 \rangle$
- v is protected by  $\{l_1\}$ .

- $$\begin{split} \mathsf{T}_{1}: \; & \mathsf{lock}(\&l_{1}); \\ & \nu = \nu + 1; \\ & \mathsf{unlock}(\&l_{1}); \\ & \mathsf{T}_{2}:\; & \mathsf{lock}(\&l_{2}); \\ & \nu = \nu + 1; \end{split}$$
  - v v + 1, unlock(& $l_2$ );
- List of accesses:  $\langle \nu, \{l_1\}, write, file.c: 2 \rangle$  $\langle \nu, \{l_2\}, write, file.c: 5 \rangle$
- No common lock!



## **Goblint: Race Detection World Champion** Potentially somewhat useful as well ...

#### Targeted Static Analysis for OCaml C Stubs: Eliminating gremlins from the code

OCaml 2023
Sat 9 Sep 2023 10:07 - 10:30 at Grand Crescent - Session 1 Chair(s): Benoît Montagu
Migration to OCaml 5 requires updating a lot of C bindings due to the removal of naked pointer support. Writing OCaml user-defined primitives in C is a necessity, but is unsafe and error-prone. It does not benefit from either OCaml's or C's type checking, and existing C static analysers are not aware of the OCaml GC safety rules, and cannot infer them from existing macros alone. The alternative is automatically generating C stubs, which requires correctly managing value lifetimes. Having a static analyser for OCaml to C interfaces is useful outside the OCaml 5 porting effort too.
After some motivating examples of real bugs in C bindings a static analyser is presented that finds these known classes of bugs. The tool works on the OCaml abstract parse and typed trees, and generates a header file and a caller model. Together with a simplified model of the OCaml runtime this is used as input to a static analysis framework, Goblint. An analysis is developed that tracks dereferences of OCaml values, and together with the existing framework reports incorrect dereferences. An example is shown how to extend the analysis to cover more safety properties.

The tools and runtime models are generic and could be reused with other static analysis tools.



- 1. Goblint (1304)
- 2. Deagle (1211)
- 3. Dartagnan (768)
- 4. UAutomizer (756)
- 5. UGemCutter (732)
- 6. UTaipan (612)
- 7. CPAchecker (400)
- 8. Locksmith (226)
- 9. Theta (205)
- 10. ...



# Why just "somewhat useful"

- When running our SV-COMP configuration on top Github repos.
- Inconclusive results in 2h.
- One clearly needs precision adjustment...
- And targeted abstractions!

Goblint 2023-11-10 16:13:46 UTC goblint.svcomp.Concrat							
status	cputime (s)	walltime (s)	memory (MB)	safe	vulnerable	uns	
Show all ~	Min:Max	Min:Max	Min:Ma	Min	Min:Max	Mir	
TIMEOUT	7200	7200	13800				
unknown	16.3	16.3	49.7	50	1		
TIMEOUT	7200	7200	14200				
unknown	48.6	48.6	360	29	0		
unknown	3480	3480	5730	118	0		
TIMEOUT	7200	7200	9950				
TIMEOUT	7200	7200	979				
true	.186	.187	26.6	2	0		
EXCEPTION (Stack overflow)	90.1	90.1	384				
EXCEPTION (Failure)	.173	.173	30.9				
unknown	2320	2320	11400	233	1		
EXCEPTION (Stack overflow)	1120	1120	3160				
TIMEOUT	7200	7200	21600				
EXCEPTION (Stack overflow)	939	939	3240				
EXCEPTION (Stack overflow)	34.6	34.6	191				
EXCEPTION (Stack overflow)	1120	1120	7660				
SEGMENTATION FAULT	2320	2320	3880				
true	4.86	4.86	42.4	83	0		



## Strength: Locks

- Infer the locked addresses
   locks[i]
- Information about pointers  $pos \in slot[i]$
- Disjointness information slot[i]  $\cap$  slot[j] =  $\emptyset$

es rs

locks

slots

pos

### Failings... **Non-Locking Concurrency**

- Thread-creation and joining
- Data segmentation
- When the number of threads is unbounded, this is hard...
- These "real-world" race challenges were submitted to SV-COMP!
- <u>https://github.com/goblint/bench/blob/master/</u> concrat/race-challenges/README.md

```
int *datas;
void *thread(void *arg) {
  int i = (int)arg;
  datas[i] = ...; // No locking needed
  return NULL;
int main() {
  int threads_total = __VERIFIER_nondet_int();
  pthread_t *tids = malloc(threads_total * sizeof(...));
  datas = malloc(threads_total * sizeof(int));
  // create threads
  for (int i = 0; i < threads_total; i++) {</pre>
    pthread_create(&tids[i], NULL, &thread, (void*)i);
  // join threads
  for (int i = 0; i < threads_total; i++) {</pre>
    pthread_join(tids[i], NULL);
  // compute with data - no locking needed
```

### **Towards usable analysis** Feedback from one industrial partner...

- False alarm can still be inisightful and useful part of code review, but...
- "Why do you produce so many warnings about the same issue?"
- We (and perhaps others too) have not paid that much attention to explaining verification outcomes.
- Heuristic analyzers are much nicer and give actionable feedback.

## Interactive analysis

- Incremental abstract interpretation
- GUI integration via MagPie Bridge
- Server mode for Goblint
- Astronomical speedup (for superficial analysis)
- Modest speedup for SV-COMP quality analyses



## **Abstract Debugging** More familiar experince of analysis results... (??)

```
\sim watch
```

```
c.m[2]: [Not {0}([0,7]), [79,81], Z]
(c.n * 25 + 11) % 100: [Unknown int([-31,31]), [1
c.n > 0: [1, [1,1], 1]
```

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	C variables_example.c > 分 main()				
+7Z]} Z], length: [3, […	2				
	3	<pre>struct s {</pre>			
	4	int n;			
	5	<pre>int m[3];</pre>			
	6	};			
	7				
		show arg   show cfg			
	8	int main() {			
	9	<pre>int a = rand() % 19;</pre>			
	10	int b = 79;			
	11	<pre>struct s c = {</pre>			
	12	7 * a + 5,			
	13	{b, b + 1, b + 2}			
	14	};			
	15	Dreturn 0;			
11,86], 11+25Z]	16	}			